

# Local Time of Transient Random Walk on $\mathbb{Z}_+$

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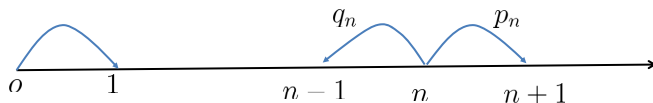
# The Model

## ◇ (1,1) Spatially Inhomogeneous Random Walks

$$P(X_{k+1} = 1 | X_k = 0) = 1,$$

$$P(X_{k+1} = n + 1 | X_k = n) = p_n,$$

$$P(X_{k+1} = n - 1 | X_k = n) = q_n,$$



where  $\forall n \geq 1, p_n, q_n > 0, p_n + q_n = 1$ .

### Transience criterion

Let  $\rho_n = q_n/p_n, n \geq 1$ . Then the chain  $X$  is **transient** if and only if  $\sum_{k=1}^{\infty} \rho_1 \cdots \rho_k < \infty$ .

# Cutpoints and Strong Cutpoints

## ◇ Definition(**Local Time**)

For  $x \in \mathbb{Z}_+$ , we call  $\xi(x) = \sum_{k=0}^{\infty} 1_{\{X_k=x\}}$  the local time of the chain  $X$  at  $x$ .

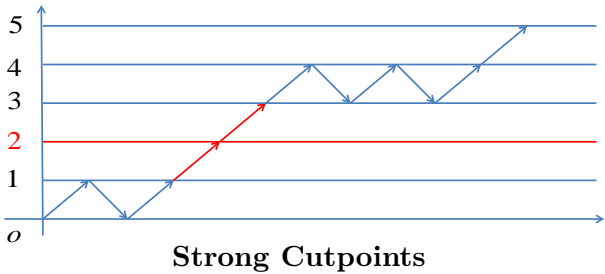
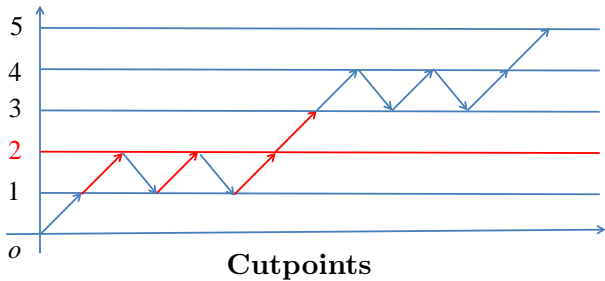
## ◇ Definition(**Cutpoint**)

A site  $R$  is a cutpoint if, for some  $k$ , we have  $X_k = R$ , and  $\{X_0, X_1, \dots, X_k\}$  is disjoint from  $\{X_{k+1}, X_{k+2}, \dots\}$ , i.e.,  $X_i \leq R, i = 0, 1, \dots, k$ ;  $X_k = R$ ; and  $X_i > R, i = k + 1, k + 2, \dots$

## ◇ Definition(**Strong Cutpoint**)

A site  $R$  is a strong cutpoint if, for some  $k$ , we have  $X_k = R, X_i < R, i = 0, 1, \dots, k - 1$ , and  $X_i > R, i = k + 1, k + 2, \dots$

♡ *Remark.* A site  $R$  is a strong cutpoint iff  $\xi(R) = 1$ .



- ♡ For recurrent random walk, there is no cutpoint.
- ♡ For transient simple(spatially homogeneous) random walk, there must be infinitely many cutpoints.
- ♡ Intuitively, the **faster** the walk runs, the **more cutpoints** it has.
- ♡ Is it possible that the walk is **transient**, but not fast enough, so that there are only **finitely many cutpoints**?

- ◇ For the above model, James, Lyons, Peres (2008) give an example which shows that the walk is **transient** but has only **finitely** many cutpoints.
- ◇ Csáki, Földes, Révész (2010) give a **criterion** for the finiteness of the number of cutpoints.



E. Csáki, A. Földes, P. Révész, On the number of cutpoints of transient nearest neighbor random walk on the line, *J. Theoret. Probab.* 23 (2) (2010) 624-638.



N. James, R. Lyons, Y. Peres, A transient Markov chain with finitely many cutpoints, *In: IMS Collections Probability and Statistics: Essays in Honor of David A. Freedman*, 2 (2008) 24-29, Institute of Mathematical Statistics.

Set for  $n \in \mathbb{Z}_+$ ,

$$\rho_n = \frac{q_n}{p_n} \text{ and } D(n) = 1 + \sum_{j=1}^{\infty} \rho_{n+1} \cdots \rho_{n+j}.$$

### Theorem(Csáki, Földes, Révész (2010))

Suppose  $0 \leq p_i < 1/2, i \geq 1$ .

- If

$$\sum_{n=1}^{\infty} \frac{1}{D(n) \log n} < \infty,$$

then almost surely,  $\{X_n\}$  has **finitely** many cutpoints;

- If  $\exists \delta > 0$  such that  $D(n) \leq \delta n \log n$  for  $n$  large enough and

$$\sum_{n=1}^{\infty} \frac{1}{D(n) \log n} = \infty,$$

then almost surely,  $\{X_n\}$  has **infinitely** many strong cutpoints.



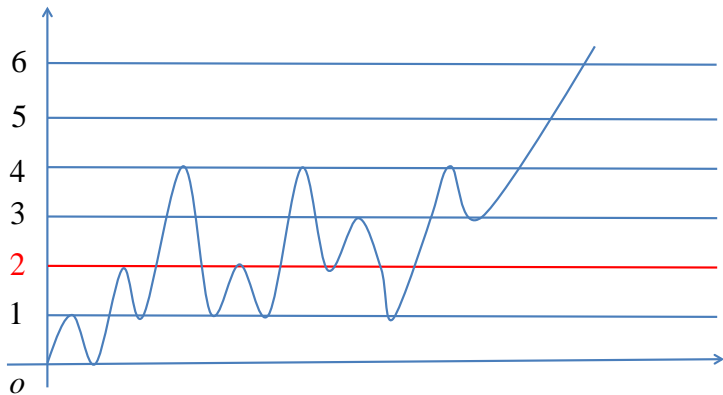
## Some Questions

Csáki, Földes, Révész proposed the following open problem:

- ◇ **Open problem**(Csáki, Földes, Révész (2010)). Consider a spatially inhomogeneous random walk. For  $x \in \mathbb{Z}_+$ , let  $\xi(x)$  be the local time of the walk at  $x$ . For any positive integer  $a$ , is there a criterion to determine whether the cardinality of the set  $C := \{x \in \mathbb{Z}^+ : \xi(x) = a\}$  is finite or infinity?

Beside the above open problem, naturally, one may ask the following question:

- ◇ **Question.** Whenever there are infinitely many (strong) cutpoints, how many cutpoints are there in  $[0, n]$ ?



$a = 8$  and  $\xi(2) = a$ ,  
the walk hits the site 2 exactly 8 times.

# The Main Results

For any  $a \in \mathbb{Z}_+$ , set  $C := \{x \in \mathbb{Z}^+ : \xi(x) = a\}$ .

## Theorem 1

Suppose that  $\rho_k$  is increasing in  $k > N_0$  for some  $N_0 > 0$  and  $\rho_k \rightarrow 1$  as  $k \rightarrow \infty$ . If

$$\sum_{n=2}^{\infty} \frac{1}{D(n) \log n} < \infty,$$

then  $|C| < \infty$  almost surely; if  $D(n) \leq \delta n \log n$ ,  $n \geq n_0$  for some  $n_0 > 0$  and  $\delta > 0$  and

$$\sum_{n=2}^{\infty} \frac{1}{D(n) \log n} = \infty,$$

then  $|C| = \infty$  almost surely.

Theorem 1 gives an answer to the above mentioned open problem proposed by Csáki, Földes, Révész (2010).

## A Concrete Example

Fix  $\beta \in \mathbb{R}$ . Clearly there exists  $n_1 > 0$  such that  $\frac{1}{4} \left( \frac{1}{n} + \frac{1}{n(\log \log n)^\beta} \right) \in (0, 1/2)$  for all  $n > n_1$ . For  $n \geq 1$  set

$$r_n = \begin{cases} \frac{1}{4} \left( \frac{1}{n} + \frac{1}{n(\log \log n)^\beta} \right), & \text{if } n \geq n_1, \\ r_{n_1}, & \text{if } n < n_1. \end{cases} \quad (1)$$

### Lemma 1

For  $n \geq 1$  set  $p_n = \frac{1}{2} + r_n$ . Then  $D(n) \sim n(\log \log n)^\beta$  as  $n \rightarrow \infty$ .

### Corollary 1

For  $n \geq 1$  set  $p_n = \frac{1}{2} + r_n$ . Then

$\beta > 1 \Rightarrow |C| < \infty$  almost surely;

$\beta \leq 1 \Rightarrow |C| = \infty$  almost surely.

## Number of Strong Cutpoints in $[0, n]$

### Proposition 1

For  $n \geq 1$  let  $r_n$  be the one in (1) and set  $p_n = \frac{1}{2} + r_n$ . Then

$$\lim_{n \rightarrow \infty} \frac{2(\log \log n)^\beta}{\log n} E |C \cap [0, n]| = 1.$$

Inspired by Proposition 1, one may expect that

$$\frac{2(\log \log n)^\beta}{\log n} |C \cap [0, n]| \rightarrow 1$$

almost surely as  $n \rightarrow \infty$ . However this is not the case. The main reason is that the events  $\{x \in C\}, x = 1, 2, \dots, n$  are not independent. We have the following theorem.

## Theorem 2

Suppose that  $a = 1$ . Let  $r_1 = 1/4$  and  $r_n = \frac{1}{2n}, n \geq 2$ . Set  $p_n = \frac{1}{2} + r_n$ , for  $n \geq 1$ . Then

$$\lim_{n \rightarrow \infty} \frac{2|C \cap [0, n]|}{\log n} \stackrel{\mathcal{D}}{=} S$$

where  $S$  is an exponentially distributed random variable with  $P(S > t) = e^{-t}, t > 0$ .

As remarked above, when  $a = 1$ , the set  $C$  is the collection of strong cutpoints. For cutpoints defined above, similar result holds. Let  $\tilde{C}$  be the collection of **cutpoints**. Under the conditions of Theorem 2, we have

$$\lim_{n \rightarrow \infty} \frac{|\tilde{C} \cap [0, n]|}{\log n} \stackrel{\mathcal{D}}{=} S.$$

# Ideas of Proofs

## Idea of the Proof of Theorem 1

For any  $a \in \mathbb{Z}_+$ , set  $C := \{x \in \mathbb{Z}^+ : \xi(x) = a\}$ .

### Theorem 1

Suppose that  $\rho_k$  is increasing in  $k > N_0$  for some  $N_0 > 0$  and  $\rho_k \rightarrow 1$  as  $k \rightarrow \infty$ . If

$$\sum_{n=2}^{\infty} \frac{1}{D(n) \log n} < \infty,$$

then  $|C| < \infty$  almost surely; if  $D(n) \leq \delta n \log n$ ,  $n \geq n_0$  for some  $n_0 > 0$  and  $\delta > 0$  and

$$\sum_{n=2}^{\infty} \frac{1}{D(n) \log n} = \infty,$$

then  $|C| = \infty$  almost surely.

### Step 1: Distribution of Local Times

For  $n > m \geq 0$ , write

$$D(m, n) := 1 + \sum_{j=1}^{n-m-1} \rho_{m+1} \cdots \rho_{m+j}, \quad D(m) := \lim_{n \rightarrow \infty} D(m, n).$$

#### Lemma 2

We have for  $x \geq 0$  and  $a \geq 1$ ,

$$P(\xi(x) = a) = \frac{p_x}{D(x)} \left( 1 - \frac{p_x}{D(x)} \right)^{a-1},$$

and for  $m \geq 2$ ,  $1 \leq j_1 < j_2 < \dots < j_m$ ,

$$\begin{aligned} P(\xi(j_1) = 1, \dots, \xi(j_m) = 1) \\ = \frac{p_{j_1} p_{j_2} \cdots p_{j_m}}{D(j_1, j_2) D(j_2, j_3) \cdots D(j_{m-1}, j_m) D(j_m)}. \end{aligned}$$



*Step 2: Cardinalities of Certain Simplexes in  $\mathbb{Z}_+^j$*

For positive integers  $a \geq j \geq i \geq 1$  set

$$S(a, j) = \left\{ (a_1, \dots, a_j) : a_k \in \mathbb{Z}_+ / \{0\}, j \geq k \geq 1, \sum_{k=1}^j a_k = a \right\},$$

$$\tilde{S}(a, j) = \left\{ (a_1, \dots, a_j) : a_k \in \mathbb{Z}_+, j \geq k \geq 1, a_j \geq 1, \sum_{k=1}^j a_k = a \right\},$$

$$\tilde{S}_i(a, j) = \left\{ (a_1, \dots, a_j) \in \tilde{S}(a, j) : \sum_{k=1}^j 1_{a_k \neq 0} = i \right\}.$$

### Lemma 3

For  $a \geq j \geq i \geq 1$ , we have  $|S(a, j)| = \binom{a-1}{j-1}$  and  $|\tilde{S}_i(a, j)| = \binom{j-1}{i-1} \binom{a-1}{i-1}$ .

*Step 3: Joint Probability of  $\{x \in C\}$  and  $\{y \in C\}$*

Based on Lemma 2 and Lemma 3, we have

**Lemma 4**

For  $1 \leq x < y < \infty$  and  $a \geq 1$ , we have

$$\begin{aligned} &P(\xi(x) = a, \xi(y) = a) \\ &= \sum_{i=1}^a \binom{a-1}{i-1} \left(\frac{p_x}{D(x,y)}\right)^i \left(1 - \frac{p_x}{D(x,y)}\right)^{a-i} \\ &\quad \times \left(q_y \left(1 - \frac{D(x,y-1)}{D(x,y)}\right)\right)^{i-1} \\ &\quad \times \left(q_y \frac{D(x,y-1)}{D(x,y)} + p_y \left(1 - \frac{1}{D(y)}\right)\right)^{a-i} \frac{p_y}{D(y)}. \end{aligned}$$

*Step 4: Dependence of  $\{x \in C\}$  and  $\{y \in C\}$*

Based on Lemma 4, we can show

**Proposition 2**

For each  $\varepsilon > 0$ , there exist  $N > 0$  and  $M > 0$  such that

$$1 - \varepsilon \leq \frac{D(x, x + y)}{D(x)} \frac{P(\xi(x) = a, \xi_{x+y} = a)}{P(\xi(x) = a)P(\xi(x + y) = a)} \leq 1 + \varepsilon,$$

for all  $x > N, y > M$ .

Based on Lemma 2 and Proposition 2, roughly speaking, Theorem 1 can be proved by **Borel-Cantelli** lemma.  $\square$

## Idea of the Proof of Theorem 2

### Theorem 2

Suppose that  $a = 1$ . Let  $r_1 = 1/4$  and  $r_n = \frac{1}{2n}, n \geq 2$ . Set  $p_n = \frac{1}{2} + r_n$ , for  $n \geq 1$ . Then

$$\lim_{n \rightarrow \infty} \frac{2|C \cap [0, n]|}{\log n} \stackrel{D}{=} S$$

where  $S$  is an exponentially distributed random variable with  $P(S > t) = e^{-t}, t > 0$ .

Theorem 2 is proved by an **moment method**. The key step is to show that the **moments of  $2|C \cap [0, n]|$**  converge to **those of exponential distribution**, which is very complicated even if we consider here only the case  $a = 1$  and  $\beta = 0$ .

## Convergence of the Moments

### Lemma 5

Under the conditions of Theorem 2, we have

$$\lim_{n \rightarrow \infty} \frac{E(|C \cap [0, n]|^k)}{(\log n)^k} = \frac{2^k}{k!}, \quad k \geq 1.$$

### Sketch of Proof.

For  $k \geq 1$ , we set  $\eta_k = \begin{cases} 1, & k \in C \\ 0 & k \notin C \end{cases}$ . Then  $|C \cap [0, n]| = \sum_{k=0}^n \eta_k$ .

Let  $S(a, j)$  be the above defined simplex in  $\mathbb{Z}_+^j$ . Then

$$\begin{aligned}
E|C \cap [0, n]|^k &= E \left( \left( \sum_{j=0}^n \eta_j \right)^k \right) \\
&= \sum_{1 \leq j_1, j_2, \dots, j_k \leq n} E(\eta_{j_1} \eta_{j_2} \cdots \eta_{j_k}) \\
&= \sum_{m=1}^k \sum_{\substack{l_1 + \dots + l_m = k, \\ l_i \geq 1, i = 1, \dots, m}} \sum_{0 \leq j_1 < \dots < j_m \leq n} m! E(\eta_{j_1}^{l_1} \cdots \eta_{j_m}^{l_m}) \\
&= \sum_{m=1}^k \sum_{(l_1, \dots, l_m) \in \mathcal{S}(k, m)} m! \sum_{0 \leq j_1 < \dots < j_m \leq n} E(\eta_{j_1}^{l_1} \cdots \eta_{j_m}^{l_m}).
\end{aligned}$$

Using the above Lemma 2([local time distribution](#)) and Lemma 3([Combinatorial result](#)), we get

$$\begin{aligned}
E|C \cap [0, n]|^k &= \sum_{m=1}^k \binom{k-1}{m-1} m! \sum_{0 \leq j_1 < \dots < j_m \leq n} E(\eta_{j_1} \cdots \eta_{j_m}) \\
&= \sum_{m=1}^k \binom{k-1}{m-1} m! \sum_{0 \leq j_1 < \dots < j_m \leq n} \frac{p_{j_1} p_{j_2} \cdots p_{j_m}}{\prod_{s=1}^{m-1} D(j_s, j_{s+1}) D(j_m)} \\
&=: \sum_{m=1}^k \binom{k-1}{m-1} m! G(n, m).
\end{aligned}$$

What is left for us to compute is

$$G(n, m) := \sum_{0 \leq j_1 < \dots < j_m \leq n} \frac{p_{j_1} p_{j_2} \cdots p_{j_m}}{\prod_{s=1}^{m-1} D(j_s, j_{s+1}) D(j_m)}.$$

We need the follow estimations of  $D(i)$  and  $D(i, j)$ , and another combinatorial result.

## Lemma 6

Under the conditions of Theorem 2, there exists a number  $i_0 > 0$  such that

$$(1 - \varepsilon) \frac{i(j-i)}{j} \leq D(i, j) \leq (1 + \varepsilon) \frac{i(j-i)}{j}, \quad j > i \geq i_0,$$
$$(1 - \varepsilon)i \leq D(i) \leq (1 + \varepsilon)i, \quad i \geq i_0.$$

## Lemma 7

Fix  $k \geq 1, l \geq 1$ . We have

$$\lim_{n \rightarrow \infty} \frac{1}{(\log n)^k} \sum_{l \leq j_1 < \dots < j_k \leq n} \frac{1}{j_1(j_2 - j_1) \cdots (j_k - j_{k-1})} = 1.$$



With the help of Lemma 6 and Lemma 7, we can show






$$\lim_{n \rightarrow \infty} \frac{G(n, m)}{(\log n)^m} = \frac{1}{2^m}$$

which leads to






$$\lim_{n \rightarrow \infty} \frac{E|C \cap [0, n]|^k}{(\log n)^k} = \frac{k!}{2^k}.$$

□

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